



Research Paper

ON THE SUBSPACE DISTANCE OF THE SUBSPACE CODES

SEYEDEH HAWRA SADROLHOFFAZ AND REZA KAHKESHANI\*

ABSTRACT. Let  $\mathcal{P}_q(n)$  be the set of all subspaces in the vector space  $\mathbb{F}_q^n$ . There is a subspace distance  $d_S(U, V)$  between any two subspaces  $U$  and  $V$ . A subspace code is also a subset of  $\mathcal{P}_q(n)$ . It is known that  $d_S(U, V) \geq d_H(\nu(\pi U), \nu(\pi V))$ , where  $\pi \in S_n$ ,  $\nu(U)$  denotes the pivot vector of  $E(U)$  and  $E(U)$  is the reduced row echelon form of the generator matrix of  $U$ . In this paper, we show that if  $E(U)$  and  $E(V)$  have at most one non-zero entry in each rows and each columns then the equality holds. Moreover, we introduce the sets  $\mathcal{G}_{U,V} = \{\pi \in S_n \mid d_S(U, V) = d_H(\nu(\pi U), \nu(\pi V))\}$  for any  $U, V \in \mathcal{P}_q(n)$  and examine them in the spaces  $\mathcal{P}_2(4)$ ,  $\mathcal{P}_2(5)$ ,  $\mathcal{P}_2(6)$  and  $\mathcal{P}_3(4)$ . It is shown that the groups  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, S_4$  and  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, D_8, S_3 \times \mathbb{Z}_2, S_4, S_5$  appears between these sets in  $\mathcal{P}_2(4)$  and  $\mathcal{P}_2(5)$ , respectively. Moreover, the groups  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, D_8, \mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_2, S_3 \times \mathbb{Z}_2, D_8 \times \mathbb{Z}_2, S_4, S_3 \times S_3, S_4 \times \mathbb{Z}_2, (S_3 \times S_3):2, S_5, S_6$  and  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, D_8, S_4$  appears between these sets in  $\mathcal{P}_2(6)$  and  $\mathcal{P}_3(4)$ , respectively.

DOI: 10.22034/as.2024.20208.1646

MSC(2010): Primary: 94B60, Secondary: 05A05, 05B20.

Keywords: Pivot vector, Subspace code, Subspace distance.

Received: 16 June 2023, Accepted: 01 September 2024.

\*Corresponding author

## 1. INTRODUCTION

Let  $\Sigma$  be a non-empty set and  $n \geq 1$  be an integer. Any subset  $\mathcal{C}$  of  $\Sigma^n$  is called a block code of length  $n$  over the alphabet  $\Sigma$ . The elements of  $\mathcal{C}$  are called codewords. The Hamming distance between two words  $\mathbf{w} = w_1w_2 \cdots w_n$  and  $\mathbf{w}' = w'_1w'_2 \cdots w'_n$  in  $\Sigma^n$  is

$$d_H(\mathbf{w}, \mathbf{w}') = \#\{i \mid 1 \leq i \leq n, w_i \neq w'_i\},$$

and the minimum (Hamming) distance of  $\mathcal{C}$  is given by

$$d_H(\mathcal{C}) = \min\{d_H(\mathbf{c}, \mathbf{c}') \mid \mathbf{c}, \mathbf{c}' \in \mathcal{C}, \mathbf{c} \neq \mathbf{c}'\}.$$

When the alphabet is the finite field  $\mathbb{F}_q$  of order  $q$ , the code  $\mathcal{C}$  is called linear if it is closed under addition and scalar multiplication. The code  $\mathcal{C}$  is binary if  $\Sigma = \mathbb{F}_2$ . The (Hamming) weight of the word  $\mathbf{w} = w_1w_2 \cdots w_n$  is defined by  $\text{wt}(\mathbf{w}) = \#\{i \mid 1 \leq i \leq n, w_i \neq 0\}$ . Clearly,  $\text{wt}(\mathbf{w}) = d_H(\mathbf{w}, \mathbf{0})$ . Also, set  $\mathbf{w} * \mathbf{w}' = (w_1w'_1)(w_2w'_2) \cdots (w_nw'_n)$ , where  $\mathbf{w} = w_1w_2 \cdots w_n$  and  $\mathbf{w}' = w'_1w'_2 \cdots w'_n$  are two words in  $\mathbb{F}_q^n$ . See [4, 5] for more details.

Let  $\mathbb{F}_q^n$  be the vector space of dimension  $n \geq 0$  over  $\mathbb{F}_q$ . We denote by  $\mathcal{P}_q(n)$  the set of all subspaces in  $\mathbb{F}_q^n$ . Each subset of  $\mathcal{P}_q(n)$  is called a subspace code. See [2, 6] for some applications of subspace codes in random linear network coding, cryptography and distributed storage. The subspace distance between two subspaces  $U, V \in \mathcal{P}_q(n)$  is defined by

$$\begin{aligned} d_S(U, V) &= \dim(U + V) - \dim(U \cap V) \\ &= \dim(U) + \dim(V) - 2 \dim(U \cap V). \end{aligned}$$

If  $\mathcal{C}$  is a subspace code then the minimum subspace distance of  $\mathcal{C}$  is

$$d_S(\mathcal{C}) = \min\{d_S(U, V) \mid U, V \in \mathcal{C}, U \neq V\}.$$

An  $k \times n$  matrix  $G$  over  $\mathbb{F}_q$  is called a generator matrix of the  $k$ -dimensional subspace  $U \in \mathcal{P}_q(n)$  if the rows of  $G$  form a basis of  $U$ , i.e.,  $U = \langle G \rangle$ . Applying the elementary row operations on  $G$ , we can obtain the generator matrix  $E(G)$  in reduced row echelon form such that  $U = \langle E(G) \rangle$ . Since such a matrix  $E(G)$  is unique for any given subspace  $U$ , we also denote it by  $E(U)$ . Denote by  $\nu(U)$  the characteristic vector of the pivot columns in  $E(U)$ . This binary vector is called a pivot vector. Clearly,  $\nu(U) \in \mathbb{F}_2^n$  and  $\nu(U)$  depends on the ordering of the positions. Note that if  $U$  is a  $k$ -dimensional subspace of  $\mathbb{F}_q^n$  then  $\text{wt}(\nu(U)) = k$ . For example, if

$$U = \left\langle \begin{pmatrix} 1 & 1 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix} \right\rangle \subseteq \mathbb{F}_3^5,$$

then

$$E(U) = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 2 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix},$$

and  $\nu(U) = (11010) \in \mathbb{F}_2^5$ . See [3] for more information. Moreover, see [7, 8] for some recent research on the subspace codes.

Let  $U$  and  $V$  be two subspaces in  $\mathcal{P}_q(n)$  with the pivot vectors  $\nu(U)$  and  $\nu(V)$ , respectively. It is known that  $d_S(U, V) \geq d_H(\nu(U), \nu(V))$  [3]. Let  $S_n$  be the symmetric group on  $\{1, \dots, n\}$ . For a matrix  $M$  of size  $k \times n$  over  $\mathbb{F}_q$  and a permutation  $\pi \in S_n$ , let  $\pi M$  denotes the matrix arising from  $M$  by permuting its columns according to  $\pi$ . If  $U$  is a subspace of  $\mathcal{P}_q(n)$  then  $\pi U$  is defined to be  $\langle \pi E(U) \rangle$ . It is easy to see that  $\langle E(U \cap V) \rangle = \langle E(U) \rangle \cap \langle E(V) \rangle$ ,  $\langle \pi E(U \cap V) \rangle = \langle \pi E(U) \rangle \cap \langle \pi E(V) \rangle$  and  $d_S(U, V) = d_S(\pi U, \pi V)$  for all  $U, V \in \mathcal{P}_q(n)$  and  $\pi \in S_n$ . Moreover,  $d_S(U, V) \geq d_H(\nu(\pi U), \nu(\pi V))$ . In this paper, we show that if  $E(U)$  and  $E(V)$  are two matrices with at most one non-zero entry in each rows and each columns then  $d_S(U, V)$  and  $d_H(\nu(\pi U), \nu(\pi V))$  are equal to each other. Moreover, we define the subset  $\mathcal{G}_{U,V} = \{\pi \in S_n \mid d_S(U, V) = d_H(\nu(\pi U), \nu(\pi V))\}$  of  $S_n$ . This set consists of permutations that achieve equality in the mentioned inequality. Generally, this set is not a group. We examine these sets for any unordered pair  $\{U, V\}$  in  $\mathcal{P}_2(4)$ ,  $\mathcal{P}_2(5)$ ,  $\mathcal{P}_2(6)$  and  $\mathcal{P}_3(4)$ . We prove that the groups  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, S_4$  and  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, D_8, S_3 \times \mathbb{Z}_2, S_4, S_5$  appears between these sets in  $\mathcal{P}_2(4)$  and  $\mathcal{P}_2(5)$ , respectively. Furthermore, the groups  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, D_8, \mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_2, S_3 \times \mathbb{Z}_2, D_8 \times \mathbb{Z}_2, S_4, S_3 \times S_3, S_4 \times \mathbb{Z}_2, (S_3 \times S_3):2, S_5, S_6$  and  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, D_8, S_4$  appears between the mentioned sets in  $\mathcal{P}_2(6)$  and  $\mathcal{P}_3(4)$ , respectively.

## 2. On the subspace distance and pivot vectors

In this section, a sufficient condition for the equality between subspace distance of two subspaces and the hamming distance of their pivot vectors is given. In fact, we have the following theorem:

**Theorem 2.1.** *Let  $U$  and  $V$  be two subspaces in  $\mathcal{P}_q(n)$ . If  $E(U)$  and  $E(V)$  are two matrices with at most one non-zero entry in each rows and each columns then  $d_S(U, V) = d_H(\nu(\pi U), \nu(\pi V))$  for all  $\pi \in S_n$ .*

*Proof.* Suppose that  $U$  and  $V$  are two subspaces of dimensions  $k$  and  $l$ , respectively. Set  $E(U) = (\alpha_{ij})_{k \times n}$ ,  $E(V) = (\beta_{hj})_{l \times n}$ ,  $\nu(U) = (a_1, a_2, \dots, a_n)$  and  $\nu(V) = (b_1, b_2, \dots, b_n)$ . Since  $E(U)$  and  $E(V)$  have at most one non-zero entry in each rows and each columns,

$$a_j = \begin{cases} 1, & \text{if } \exists 1 \leq i \leq k \alpha_{ij} \neq 0; \\ 0, & \text{otherwise;} \end{cases}$$

and

$$b_j = \begin{cases} 1, & \text{if } \exists 1 \leq h \leq l \beta_{hj} \neq 0; \\ 0, & \text{otherwise;} \end{cases}$$

for any  $1 \leq j \leq n$ . Hence,

$$\begin{aligned} \text{wt}(\nu(U) * \nu(V)) &= \text{wt}(a_1 b_1, a_2 b_2, \dots, a_n b_n) \\ &= \#\{1 \leq j \leq n \mid a_j = b_j = 1\} \\ &= \#\{1 \leq j \leq n \mid \exists 1 \leq i \leq k \exists 1 \leq h \leq l \alpha_{ij} \neq 0, \beta_{hj} \neq 0\} \\ &= \dim(\langle E(U) \rangle \cap \langle E(V) \rangle) \\ (1) \quad &= \dim(U \cap V). \end{aligned}$$

Now, let  $\pi$  be an arbitrary permutation in  $S_n$  and set  $\pi U = \pi E(U) = (\gamma_{ij})_{k \times n}$  and  $\pi V = \pi E(V) = (\delta_{hj})_{l \times n}$ . Clearly,  $\gamma_{ij} = \alpha_{i\pi^{-1}(j)}$  and  $\delta_{hj} = \beta_{h\pi^{-1}(j)}$  for all  $1 \leq i \leq k$ ,  $1 \leq h \leq l$  and  $1 \leq j \leq n$ . We can say that  $\gamma_{i\pi(j)} \neq 0$  and  $\delta_{h\pi(j)} \neq 0$  if and only if  $\alpha_{ij} \neq 0$  and  $\beta_{hj} \neq 0$ . It follows from the assumption that  $\pi E(U)$  and  $\pi E(V)$  are also matrices that they have at most one non-zero entry in each rows and each columns. By (1), we have

$$\begin{aligned} \text{wt}(\nu(\pi U) * \nu(\pi V)) &= \#\{1 \leq j \leq n \mid \exists 1 \leq i \leq k \exists 1 \leq h \leq l \gamma_{ij} \neq 0, \delta_{hj} \neq 0\} \\ &= \#\{1 \leq j \leq n \mid \exists 1 \leq i \leq k \exists 1 \leq h \leq l \alpha_{i\pi^{-1}(j)} \neq 0, \beta_{h\pi^{-1}(j)} \neq 0\} \\ &= \#\{1 \leq \pi(j) \leq n \mid \exists 1 \leq i \leq k \exists 1 \leq h \leq l \alpha_{ij} \neq 0, \beta_{hj} \neq 0\} \\ &= \text{wt}(\nu(U) * \nu(V)) \\ (2) \quad &= \dim(U \cap V). \end{aligned}$$

So, (2) implies that

$$\begin{aligned} d_H(\nu(\pi U), \nu(\pi V)) &= \text{wt}(\nu(\pi U)) + \text{wt}(\nu(\pi V)) - 2 \text{wt}(\nu(\pi U) * \nu(\pi V)) \\ &= \dim(U) + \dim(V) - 2 \dim(U \cap V) \\ &= d_S(U, V). \end{aligned}$$

Now, the proof is complete.  $\square$

The mentioned condition in Theorem 2.1 is sufficient but it is not necessary. To see this, consider the subspaces

$$U = \left\langle \begin{pmatrix} 1 & 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 \end{pmatrix} \right\rangle, \quad \text{and} \quad V = \left\langle \begin{pmatrix} 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \end{pmatrix} \right\rangle.$$

So, we have

$$E(U) = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{pmatrix}, \quad \text{and} \quad E(V) = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}.$$

This implies that  $\nu(U) = (11100)$  and  $\nu(V) = (10011)$ . By a computer program in Maple software,  $d_H(\nu(\pi U), \nu(\pi V)) = 4$  for all  $\pi \in S_5$ . On the other hand, by definition,  $U \cap V = \left\langle \begin{pmatrix} 1 & 0 & 0 & 0 & 1 \end{pmatrix} \right\rangle$  and  $d_S(U, V) = 3 + 3 - 2 \cdot 1 = 4$ . So, the equality is hold but  $E(U)$  does not satisfy in the condition of Theorem 2.1. Now, we present an example where equality does not hold.

**Example 2.2.** Consider the subspaces

$$U = \left\langle \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix} \right\rangle, \quad \text{and} \quad V = \left\langle \begin{pmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \end{pmatrix} \right\rangle.$$

So,

$$E(U) = \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{pmatrix}, \quad \text{and} \quad E(V) = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix}.$$

Clearly,  $U \cap V = \left\langle \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix} \right\rangle$  and  $\nu(U) = \nu(V) = (1110)$ . Thus,  $d_H(\nu(U), \nu(V)) = 0$  and  $d_S(U, V) = \dim(U) + \dim(V) - 2\dim(U \cap V) = 2$ . This implies that  $d_S(U, V) > d_H(\nu(U), \nu(V))$ .

### 3. ON THE SPACES $\mathcal{P}_q(n)$ FOR SOME $n$ AND $q$

Consider again the set  $\mathcal{P}_q(n)$  consisting of all subspaces of  $\mathbb{F}_q^n$ . For any  $U$  and  $V$  in  $\mathcal{P}_q(n)$ , we define

$$\mathcal{G}_{U,V} = \{\pi \in S_n \mid d_S(U, V) = d_H(\nu(\pi U), \nu(\pi V))\}.$$

The subset  $\mathcal{G}_{U,V} \subseteq S_n$  is not necessarily a group. For example, if

$$U = \left\langle \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \right\rangle, \quad \text{and} \quad V = \left\langle \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix} \right\rangle$$

then

$$\mathcal{G}_{U,V} = \{(), (12), (34), (23), (13), (24), (123), (132), (234), (243), (12)(34), (1342)\}$$

and clearly, this set is not a group. Also, consider the subspaces  $U$  and  $V$  of example 2.2. Since  $1 \notin \mathcal{G}_{U,V}$ , the set  $\mathcal{G}_{U,V}$  is not a group. Now, a consequence of Theorem 2.1 is as follows:

**Corollary 3.1.** *If  $U$  and  $V$  satisfy in the assumptions of Theorem 1 then  $\mathcal{G}_{U,V} \cong S_n$ .*

In this section, we examine these sets by a computer program in Maple software, where  $(n, q) = (4, 2), (5, 2), (6, 2), (4, 3)$ . By [1, 3],

$$|\mathcal{P}_q(n)| = \sum_{k=1}^n \begin{bmatrix} n \\ k \end{bmatrix}_q = \sum_{k=1}^n \prod_{i=0}^{k-1} \frac{q^{n-i} - 1}{q^{k-i} - 1},$$

and hence, we have  $(|\mathcal{P}_2(4)|) = 2145$ ,  $(|\mathcal{P}_2(5)|) = 69378$ ,  $(|\mathcal{P}_2(6)|) = 3986076$  and  $(|\mathcal{P}_3(4)|) = 22155$  such sets in the cases  $(n, q) = (4, 2)$ ,  $(n, q) = (5, 2)$ ,  $(n, q) = (6, 2)$  and  $(n, q) = (4, 3)$ , respectively. Among these sets, there are 602, 8030, 140912 and 2765 subgroups, respectively. The information we obtain about all such subgroups are listed in Tables 1-6. The subgroups and their numbers are written under the columns ' $\mathcal{G}_{U,V}$ ' and '#', respectively. The headings ' $\cong$ ' indicates the shape of the group  $\mathcal{G}_{U,V}$ . The last column gives an example for each case. Therefore, we obtain the following theorems:

**Theorem 3.2.** *Consider the words in the space  $\mathcal{P}_2(4)$ . There are 602 groups among the 2145 sets of  $\mathcal{G}_{U,V}$  and these groups are isomorphic to  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3$  and  $S_4$ .*

**Theorem 3.3.** *Consider the words in the space  $\mathcal{P}_2(5)$ . There are 8030 groups among the 69378 sets of  $\mathcal{G}_{U,V}$  and these groups are isomorphic to  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, D_8, S_3 \times \mathbb{Z}_2, S_4$  and  $S_5$ .*

**Theorem 3.4.** *Consider the words in the space  $\mathcal{P}_2(6)$ . There are 140912 groups among the 3986076 sets of  $\mathcal{G}_{U,V}$  and these groups are isomorphic to  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, D_8, \mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_2, S_3 \times \mathbb{Z}_2, D_8 \times \mathbb{Z}_2, S_4, S_3 \times S_3, S_4 \times \mathbb{Z}_2, (S_3 \times S_3):2, S_5$  and  $S_6$ .*

**Theorem 3.5.** *Consider the words in the space  $\mathcal{P}_3(4)$ . There are 2765 groups among the 22155 sets of  $\mathcal{G}_{U,V}$  and these groups are isomorphic to  $1, \mathbb{Z}_2, \mathbb{Z}_2 \times \mathbb{Z}_2, S_3, D_8$  and  $S_4$ .*

TABLE 1. The groups  $\mathcal{G}_{U,V}$ , where  $n = 4$  and  $q = 2$ .

$\mathcal{G}_{U,V}$	#	$\cong$	example
$\langle\langle () \rangle\rangle$	2	1	$\begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$
$\langle\langle (12) \rangle\rangle$	3	$\mathbb{Z}_2$	$\begin{pmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$
$\langle\langle (23) \rangle\rangle$	1		$\begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 \end{pmatrix}, \begin{pmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \end{pmatrix}$
$\langle\langle (34) \rangle\rangle$	3		$\begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix}, \begin{pmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$
$\langle\langle (12), (34) \rangle\rangle$	4	$\mathbb{Z}_2 \times \mathbb{Z}_2$	$\begin{pmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$
$\langle\langle (12), (123) \rangle\rangle$	6	$S_3$	$\begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$
$\langle\langle (24), (234) \rangle\rangle$	6		$\begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$
$\langle\langle (12), (1234) \rangle\rangle$	577	$S_4$	$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$

#### 4. CONCLUSION

In this paper, we show that if  $E(U)$  and  $E(V)$  have at most one non-zero entry in each rows and each columns then inequality  $d_S(U, V) \geq d_H(\nu(\pi U), \nu(\pi V))$  becomes equality. Moreover, we introduce the sets  $\mathcal{G}_{U,V} = \{\pi \in S_n \mid d_S(U, V) = d_H(\nu(\pi U), \nu(\pi V))\}$  for any  $U, V \in \mathcal{P}_q(n)$  and examine them for some  $n$  and  $q$ . Interested authors can work on finding a necessary condition for this inequality to become an equality. Also, interested authors can find the application of these results in coding theory.

#### 5. ACKNOWLEDGMENTS

The authors are grateful to the referees for their valuable comments. This work is partially supported by the University of Kashan under grant number 1073211.









TABLE 5. The groups  $\mathcal{G}_{U,V}$ , where  $n = 6$  and  $q = 2$  (Continued).

$\mathcal{G}_{U,V}$	#	$\cong$	example
$\langle(142536), (13)\rangle$	1	$(S_3 \times S_3):2$	$\left( \begin{pmatrix} 1 & 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \right)$
$\langle(12), (12345)\rangle$	158	$S_5$	$\left( \begin{pmatrix} 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \right)$
$\langle(23), (23456)\rangle$	158		
$\langle(12), (123456)\rangle$	139477	$S_6$	$\left( \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \right)$

TABLE 6. The groups  $\mathcal{G}_{U,V}$ , where  $n = 4$  and  $q = 3$ .

$\mathcal{G}_{U,V}$	#	$\cong$	example
$\langle()\rangle$	24	1	$\left( \begin{pmatrix} 1 & 2 & 0 & 2 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \right)$
$\langle(12)\rangle$	56	$\mathbb{Z}_2$	$\left( \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 2 & 0 & 2 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix} \right)$
$\langle(23)\rangle$	16		
$\langle(34)\rangle$	56		
$\langle(12), (34)\rangle$	136	$\mathbb{Z}_2 \times \mathbb{Z}_2$	$\left( \begin{pmatrix} 1 & 0 & 2 & 2 \\ 0 & 1 & 2 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix} \right)$
$\langle(12), (123)\rangle$	120	$S_3$	$\left( \begin{pmatrix} 1 & 0 & 2 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 & 0 & 2 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \right)$
$\langle(24), (234)\rangle$	120		
$\langle(34), (1324)\rangle$	8	$D_8$	$\left( \begin{pmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 2 & 2 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 2 & 0 & 1 \\ 0 & 0 & 1 & 2 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \right)$
$\langle(12), (1234)\rangle$	2229	$S_4$	$\left( \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 & 0 & 2 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \right)$

REFERENCES

[1] P. J. Cameron, *Combinatorics: Topics, Techniques, Algorithms*, Cambridge University Press, 1994.  
 [2] W. C. Huffman, J. -L. Kim and P. Solé, *Concise Encyclopedia of Coding Theory*, Chapman and Hall/CRC, 2021.  
 [3] S. Kurz, *Construction and bounds for subspace codes*, (2023) arXiv:2112.11766v2.  
 [4] S. Ling and C. Xing, *Coding Theory, A First Course*, Cambridge University Press, 2004.

- [5] F. J. MacWilliams and N. J. A. Sloane, *The Theory of Error-Correcting Codes*, North-Holland, 1977.
- [6] N. Raviv, *Subspace Codes and Distributed Storage Codes*, PhD thesis, Computer Science Department, Technion, 2017.
- [7] H. Zhang and C. Tang, *Further constructions of large cyclic subspace codes via Sidon spaces*, *Linear Algebra Appl.* **661** (2023) 106-115.
- [8] F. Zullo, *Multi-orbit cyclic subspace codes and linear sets*, *Finite Fields Their Appl.* **87** (2023) 102153.

**Seyedeh Hawra Sadrolhoffaz**

Department of Pure Mathematics, Faculty of Mathematical Sciences,  
University of Kashan  
Kashan, Iran.  
`sadr@grad.kashanu.ac.ir`

**Reza Kahkeshani**

Department of Pure Mathematics, Faculty of Mathematical Sciences,  
University of Kashan  
Kashan, Iran.  
`kahkeshanireza@kashanu.ac.ir`